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Optimising Product Swaps in Urban Retail Networks

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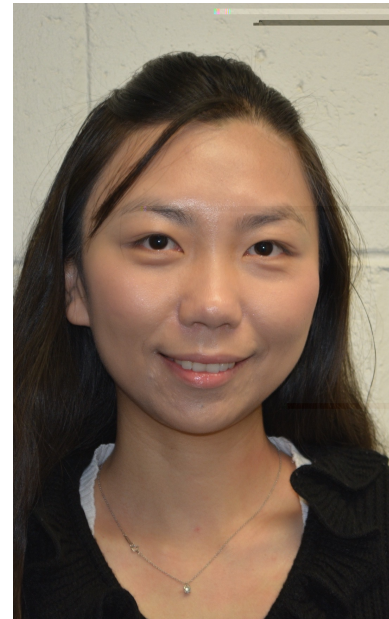
A/Prof Russell Thompson

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Dr Lele (Joyce) Zhang

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Mathematical Model

To design vehicle routes with loads and transshipment timing

Decision variables

- $y_{v,ij} \in \{0,1\} \quad v \in \mathcal{V}, i, j \in \mathcal{M}$
- $z_{v,n,ij} \in \{0,1\} \quad v \in \mathcal{V}, i, j \in \mathcal{M}, n \in \mathcal{N}$
- $\omega_{v,i} \geq 0 \quad v \in \mathcal{V}, i \in \mathcal{M}$
-

Multiple objectives

To minimise

•

$$F_{voc} = \underbrace{VC_f \sum_{v \in \mathcal{V}} \sum_{i \in \mathcal{M}} y_{v,if_v}}_{\text{red}} + \underbrace{VC_d \times V \sum_{v \in \mathcal{V}} \sum_{i \in \mathcal{M}} \sum_{j \in \mathcal{M}} y_{v,ij} T_{ij}}_{\text{blue}} \quad \left. \vphantom{F_{voc}} \right\} \$$$

•

$$F_{lc} = \underbrace{LC_d \sum_{v \in \mathcal{V}} (\overset{\downarrow}{A}_{v,f_v} - \overset{\downarrow}{\omega}_{v,f_v})}_{\text{orange}} + \underbrace{LC_f \sum_{v \in \mathcal{V}} \sum_{i \in \mathcal{M}} (\overset{\downarrow}{\mu}_{v,i} U_i + \overset{\downarrow}{\lambda}_{v,i} L_i)}_{\text{green}} \quad \left. \vphantom{F_{lc}} \right\}$$

•

$$F_{urc} = \frac{1}{2} \sum_{i \in \mathcal{M}} \sum_{v \in \mathcal{V}} \sum_{n \in \mathcal{N}} (\underbrace{p \mu_{v,n,i} + p \lambda_{v,n,i}}_{\text{purple}}) - |\mathcal{N}| \quad \swarrow$$

$$\underset{\mathbf{y}, \mathbf{z}, \boldsymbol{\omega}}{\text{Min}} (F_{\$}, F_{urc})$$



Constraints (selected, cont.)

•

$$\max_{v \in \mathcal{V}} \underbrace{\sum_{u \in \mathcal{V} \setminus \{v\}} \mathbf{1}(A_{u,i} < D_{v,i} \& A_{v,i} < D_{u,i})}_{v} + 1 \leq l_i \quad \forall i \in \mathcal{M}$$

•

$$\sum_{n \in \mathcal{N}} s_n z_{v,n,i} \leq c_v \quad \forall v \in \mathcal{V}$$

•

$$\sum_{j \in \mathcal{M} \setminus \{i\}} y_{v,i,j} = \sum_{k \in \mathcal{M} \setminus \{i\}} y_{v,k,i} \quad \forall i \in \mathcal{M}, v \in \mathcal{V}$$

$$\sum_{v \in \mathcal{V}} \sum_{j \in \mathcal{M}} z_{v,n,i,j} = \sum_{v \in \mathcal{V}} \sum_{k \in \mathcal{M}} z_{v,n,k,i} \quad \forall i \in \mathcal{M} \setminus \{o_n, d_n\}, n \in \mathcal{N}$$



Numerical Experiments

- Melbourne network with 5 nodes
- Three strategies
 - Dedicated delivery (DD)
 - No transshipment (NT)
 - Transshipment (TR)
- Solutions obtained by Gurobi 9.2



Strategy	Financial cost	# of Vehicles	VKT	Load factor

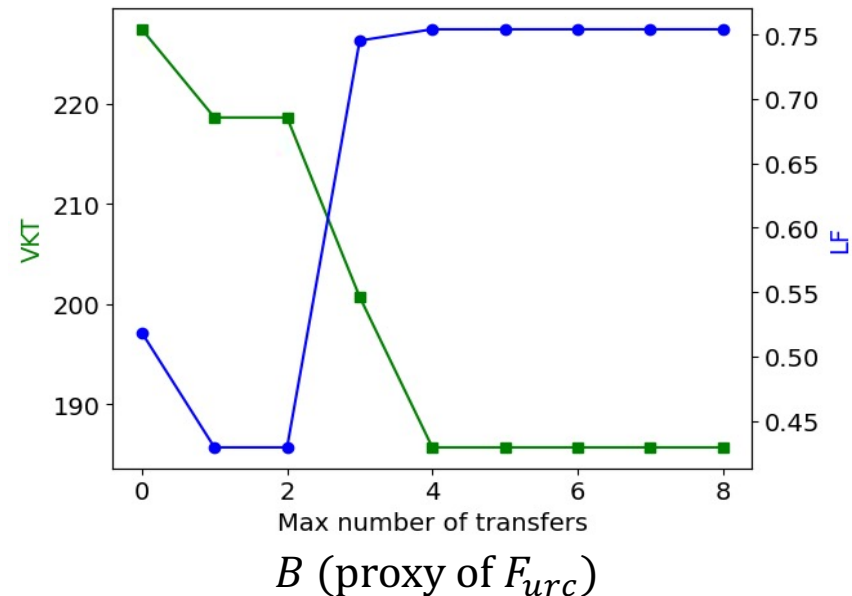
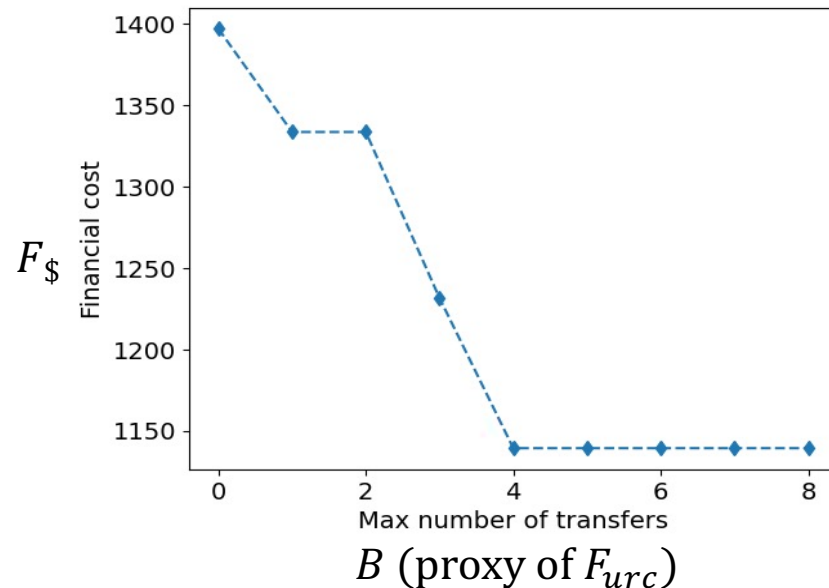


Pareto front

- Modify the multi-objective optimisation problem $\min (F_{\$}, F_{urc})$ to

$$\min F_{\$} \text{ s.t. } F_{urc} \leq B \text{ and existing constraints}$$

Upper bound of total swap number
- Financial cost and VKT decrease with B , and load factor generally increases.
- Goods transhipment benefits the delivery efficiency.





Conclusions

- Allowing products to be transhipped at stores can provide significant savings in financial cost and vehicle travel distances (and hence environmental cost).
- Designing efficient goods transfer networks requires consideration of OD patterns as well as locations of and facilities at stores.
- Transshipment may bring more uncertainties in the network. For improving the network reliability, it needs
 - Consideration of vehicle coordination;
 - Models incorporating with stochastic elements such as demand, travel time, and loading time.



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Thank you

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